

*L3 Mention Informatique
Parcours Informatique et MIAGE*

Génie Logiciel Avancé - Advanced Software Engineering

White-Box Tests (Rev)

Burkhart Wolff

(burkhart.wolff@universite-paris-saclay.fr)

<https://usr.lmf.cnrs.fr/~wolff>

Towards **Static** Program-based Unit Test

- ❑ How can we test during development
(at coding time, even at design-time ?)
- ❑ How can we test "systematically"?
 - ❑ What could be a test-generation method?
 - ❑ What could be an algorithm to generate tests?
 - ❑ What could be a coverage criterion ?
(or: adequacy criterion,
telling that we "tested enough")

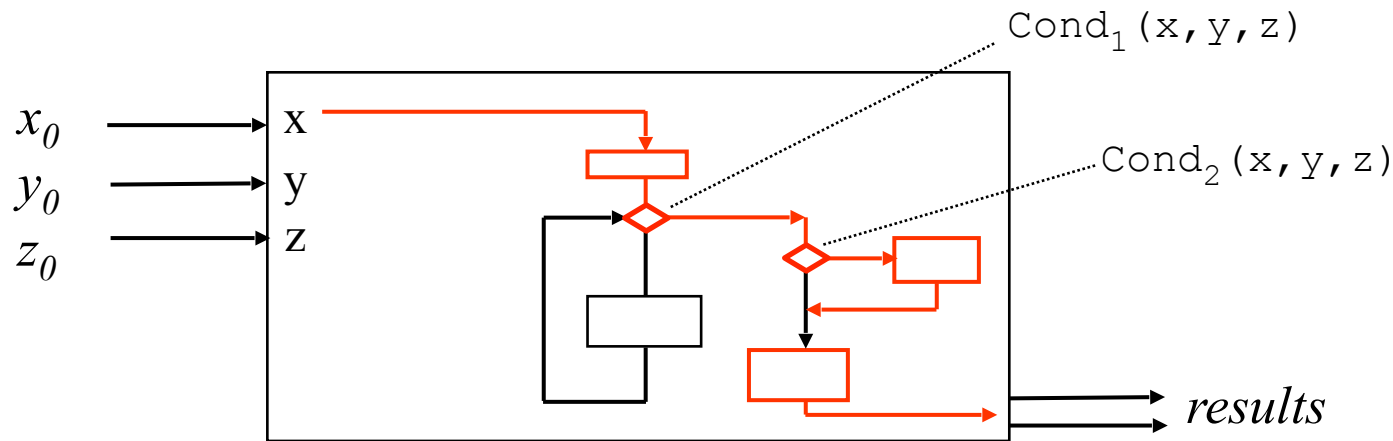
Towards **Static** Program-based Unit Test

- ❑ Let's exploit the structure of the program !!!

(and not, as before in specification based tests („black box“-tests), depend entirely on the spec).

- ❑ **Assumption:** Programmers make most likely errors in branching points of a program (Condition, While-Loop, ...), but get the program “in principle right”.
(Competent programmer assumption)
- ❑ Lets develop a test method that exploits this !

Static Structural ("white-box") Tests



□ we select "critical" paths

Idea:

□ a path corresponds to one logical expression over initial values x_0, y_0, z_0
 □ specification used to verify the obtained resultants
 corresponding to one test-case (comprising several test data)

Recall. This path-condition can be effectively constructed by a process called "symbolic execution"

We are interested either in edges (control flow), or in nodes (data flow)

The notion of a "coverage criterion"

A coverage criterion is a function mapping a CFG to a particular subset of its paths ...

- the set of paths covering all basic blocks
- the set of paths covering all conditions
- the set with all loops are traversed
- a particular subset of calls/labels occurring in the CFG has been covered
- ...

Well-known Coverage Criteria I

Criterion $C = \text{AllInstructions}(\text{CFG})$:

For all nodes N in CFG (basic instructions or decisions)
exists a path in C that contains N

Well-known Coverage Criteria II

Criterion $C = \text{AllTransitions}(\text{CFG})$:

For all arcs A in the CFG exists a path in C that uses A

Well-known Coverage Criteria III

Criterion $C = \text{AllPaths}(\text{CFG})$:

All possible paths ...

☹ Whenever there is a loop, C is infinite !

☞ weaker variant: $\text{AllPaths}_k(\text{CFG})$.

We limit the paths through a loop to maximally k times ...

☞ we have again a finite number of paths

A Hierarchy of Coverage Criteria

- $\text{AllPaths}(\text{CFG}) \supseteq$
 $\text{AllPaths}_k(\text{CFG}) \supseteq$
 $\text{AllTransitions}(\text{CFG}) \supseteq$
 $\text{AllInstructions}(\text{CFG})$
- Each of these implications reflects a proper containment; the other way round is never true.

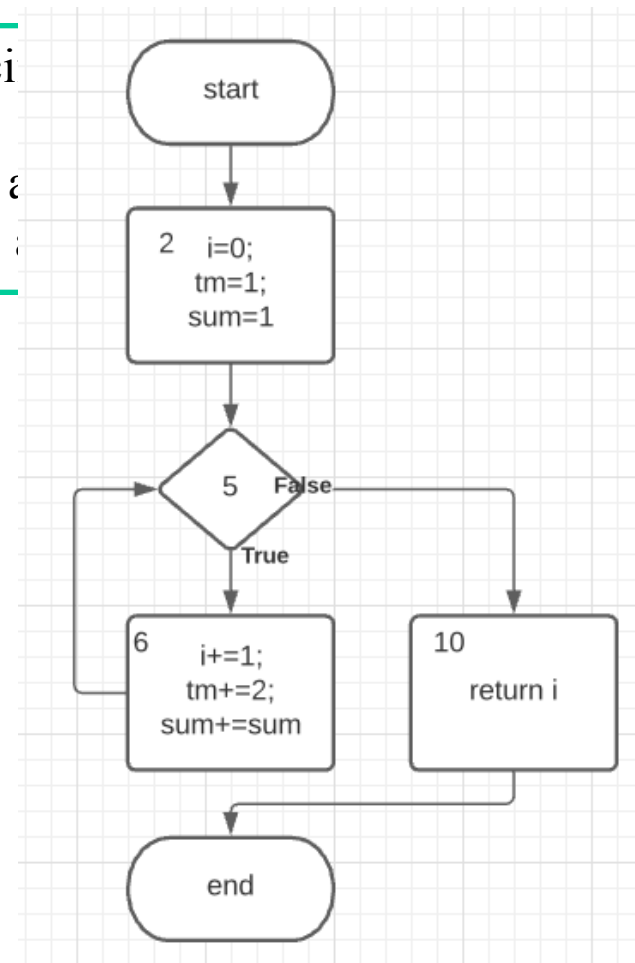
Schmankerle

□ Program:

```
int f (int a) {  
    int i = 0;  
    int tm = 1;  
    int sum = 1;  
    while(sum <= a) {  
        i = i+1;  
        tm = tm+2;  
        sum = tm+sum;  
    }  
    return i;  
}
```

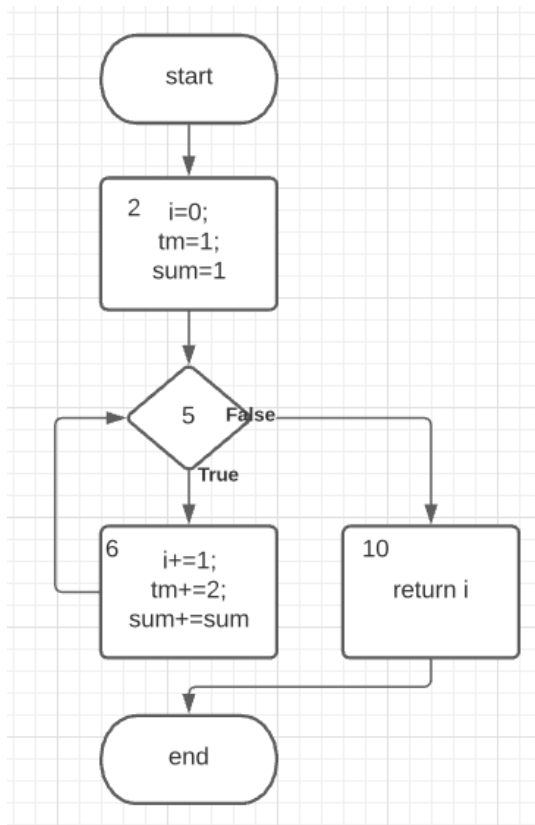
Speci

pre : a
post: i



Schmankerle

CFG de f:



For example:

$$\text{AllInstructions}(\text{CFG}) = \{\text{start}, 2, 5, 6, 5, 10, \text{end}\}$$
$$\text{AllTransitions}(\text{CFG}) = \{[\text{start}, 2, 5, 6, 5, 10, \text{end}]\}$$
$$\text{AllPath}_3(\text{CFG}) = \{[\text{start}, 2, 5, 10, \text{end}],$$
$$[\text{start}, 2, 5, 6, 5, 10, \text{end}],$$
$$[\text{start}, 2, 5, 6, 5, 6, 5, 10, \text{end}],$$
$$[\text{start}, 2, 5, 6, 5, 6, 5, 6, 5, 10, \text{end}]\}$$
$$\text{AllPath}(\text{CFG}) = \{ S \mid \exists k \in \mathbb{N} .$$

S = [start, 2, 5, (6, 5)^k, 10, end]
(infinite !)

Example: A Symbolic Path Execution

We want to execute the path from AllPath₃:

[S, 2, 5, 6, 5, 10, E]

$\Phi \mapsto$ $a_0 \geq 0$	$a_0 \geq 0$	$(\text{sum} \leq a) \sigma_2$ $\wedge a_0 \geq 0$	$1 \leq a_0 \wedge$ $a_0 \geq 0$	$1 \leq a_0 \wedge$ $\neg (\text{sum} \leq a) \sigma_6$	$1 \leq a_0 \wedge$ $4 > a_0$	$1 \leq a_0 \wedge$ $4 > a_0 \wedge$ $\text{res} = 1$
$a \mapsto a_0$	a_0	a_0	a_0	a_0	a_0	a_0
$i \mapsto i_0$	0	0	1	1	1	1
$\text{tm} \mapsto \text{tm}_0$	1	1	3	3	3	3
$\text{sum} \mapsto \text{sum}_0$	1	1	4	4	4	4

Example: A Symbolic Path Execution

Result:

Test-Case:

For the path $M=[\text{start}, 2, 5, 6, 5, 10, \text{end}]$

we have the path condition $\Phi \mapsto 1 \leq a_0 \wedge 4 > a_0$

A concrete Test, satisfying Φ :

$$\boxed{a_0 \mapsto 3}$$

Execution of program with this test vector 3: $f(3) = 1$

Verification of the post-condition: $\text{post}(3, 1) = \text{true}$

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Addendum: Multiple-Condition-Decision-Coverage

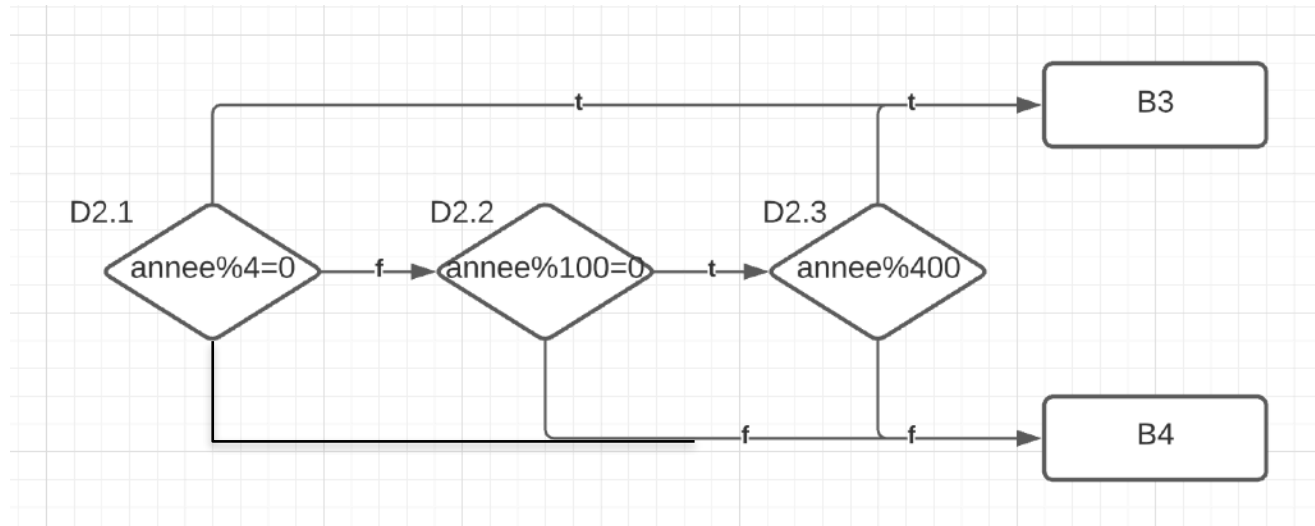
Problem: Consider:

```
if(m1 == m2) {  
    res = j2 - j1;  
} else {  
    if((annee%4 == 0) || (annee%100 == 0 && annee%400 != 0)) {  
        daysin[2] = 29;  
    } else {  
        daysin[2] = 28;  
    }  
    res = j2 + (daysin[m1] - j1);  
    for(int i = m1+1; i < m2; i++) {  
        res = res + daysin[i];  
    }  
}
```

even transition coverage on the Byzantine condition
(line 4-5) is very coarse and risks to miss the point !

Addendum: Multiple-Condition-Decision-Coverage

Solution: We use the inherent control flow in *C* for `||` and `&&` for a refined control flow graph !



now transition coverage (on the refined CFG) checks each condition individually for true and false.

This kind of coverage is called MC/DC.